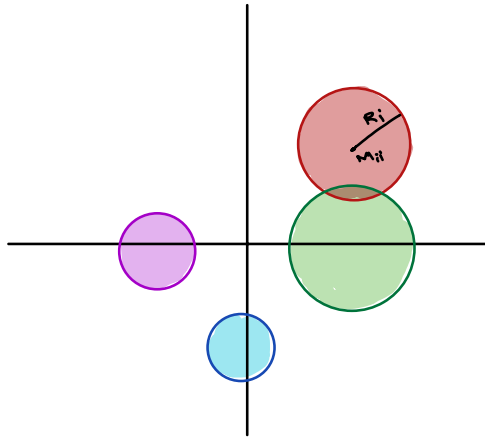


Recap

Gershgorin:



$$M \in \mathbb{C}^{n \times n}$$

eigenvalue $\lambda \in \bigcup \text{Disc}(M_{ii}, R_i)$

$$R_i = \sum_{j \neq i} |M_{ji}|$$

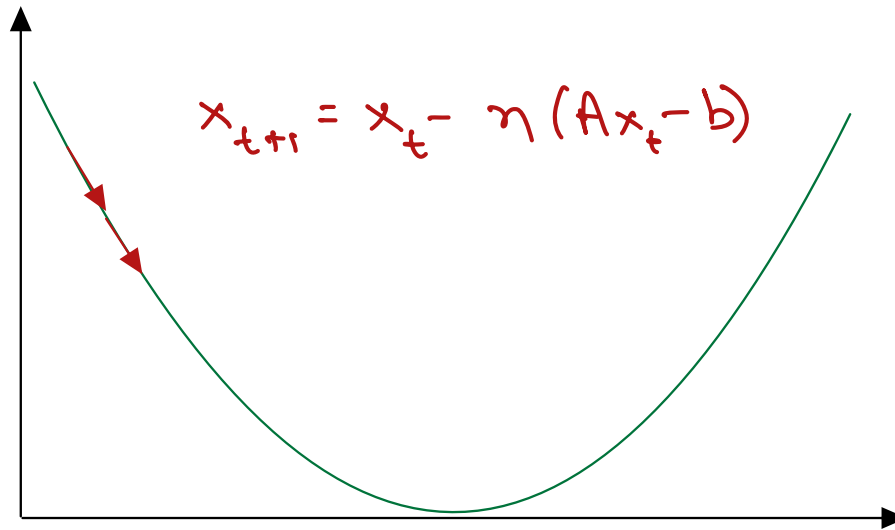
Linear Systems:

$$A_0 x = b_0$$

↓

$$A_0^T A_0 x = A_0^T b_0$$

positive definite



$$x_{t+1} = x_t - \eta(Ax_t - b)$$

$$(x_t - u) = (I - \eta A)^t (x_0 - u)$$

$$\|(I - \eta A)^t\|_2 = \left(1 - \frac{2}{\kappa + 1}\right)^t$$

$$\kappa = \frac{\lambda_1}{\lambda_n} \left\{ \begin{array}{l} \text{condition} \\ \text{number} \end{array} \right.$$

Conjugate gradient method

Steepest descent: $x_0 = 0$ (say)

$$x_1 = x_0 - \eta (Ax_0 - b) = \eta b$$

$$x_2 = x_1 - \eta (Ax_1 - b) = 2\eta b - \eta^2 Ab$$

$$x_3 = x_2 - \eta (Ax_2 - b) = 3\eta b - 3\eta^2 Ab + \eta^3 A^2 b$$

$\vdots$

$$x_t \in \text{Span}(\{b, Ab, A^2b, \dots, A^{t-1}b\})$$

Krylov space $K_t(A, b)$

What if we just found the "best vector" in $K_t(A, b)$

Minimizing in Kaylor Space

$$\text{Solve } A_0 x = b_0 \rightarrow \underbrace{A_0^T A_0}_A x = \underbrace{A_0^T b_0}_b$$

Say u is unique soln to above

$$\begin{aligned} x_t = \arg \min_{x \in K_t(A, b)} f(x) &= \arg \min_{x \in K_t(A, b)} \|A_0 x - b_0\|^2 = \langle x - u, A(x - u) \rangle \\ &= \langle x - u, x - u \rangle_A \end{aligned}$$

Let $w_1 \dots w_t$ be orthonormal basis for $K_t(A, b)$

under inner product $\langle x, y \rangle_A = \langle x, Ay \rangle$

Let optimal solution $x = \sum_{i=1}^t c_i w_i$

$$\begin{aligned} \text{Then } c_i &= \langle u, w_i \rangle_A = \langle u, Aw_i \rangle = \langle \underline{Au}, w_i \rangle \\ &= \langle b, w_i \rangle \end{aligned}$$

Convergence analysis

$$x_t = c_0 \cdot b + c_1 \cdot A b + \dots + c_{t-1} \cdot A^{t-1} b$$

$$= \underbrace{(c_0 \cdot I + c_1 \cdot A + c_2 \cdot A^2 \dots + c_{t-1} \cdot A^{t-1})}_{P(A)} \cdot \underbrace{b}_{Au}$$

say $x_0 = 0$

$$u - x_t = (I - \underbrace{P(A) \cdot A}_{q(A)}) \cdot u = (I - \underbrace{P(A) \cdot A}_{q(A)}) (u - x_0)$$

error $u - x_t = \underbrace{(I - \eta A)^t}_{\dots} (u - x_0)$

Eigenvalues

$$A : \lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_n > 0$$

$$\|q(A)\|_2$$

$$q(A): q(\lambda_1)$$

...

$$q(\lambda_n)$$

$$= \max_{i \in [n]} |q(\lambda_i)|$$

Error analysis

$$(u - x_t) = q_t(A) \cdot (u - x_0) \Rightarrow \|u - x_t\| \leq \underbrace{\|q_t(A)\|_2}_{= \max_{i \in [n]} |q_t(\lambda_i)|} \cdot \|u - x_0\|$$

Which polynomial $q_t(A)$ do we choose?

$$q_t(A) = I - A \cdot p_t(A)$$

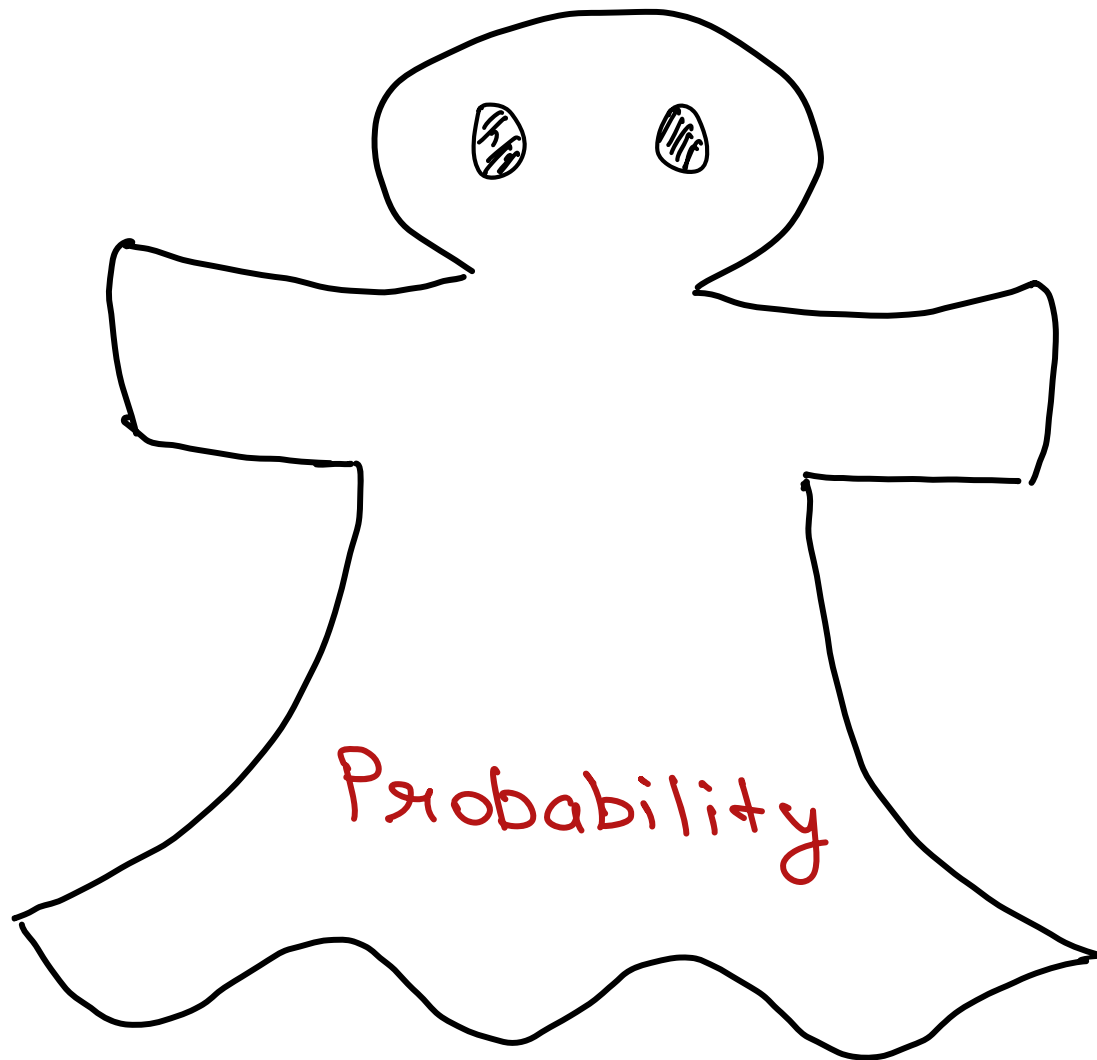
$$q_t(z) = 1 - z \cdot p_t(z)$$

► $\exists$ polynomial q , $\deg(q) \leq t$, $q(0) = 1$ s.t.

$$q(z) \leq 2 \cdot \left(1 - \frac{2}{\sqrt{\kappa} + 1}\right)^t \quad \forall z \in [\lambda_1, \lambda_n]$$

For $\|q_t(A)\|_2 \leq \epsilon$ choose $t = O(\sqrt{\kappa} \log 1/\epsilon)$

Part II :

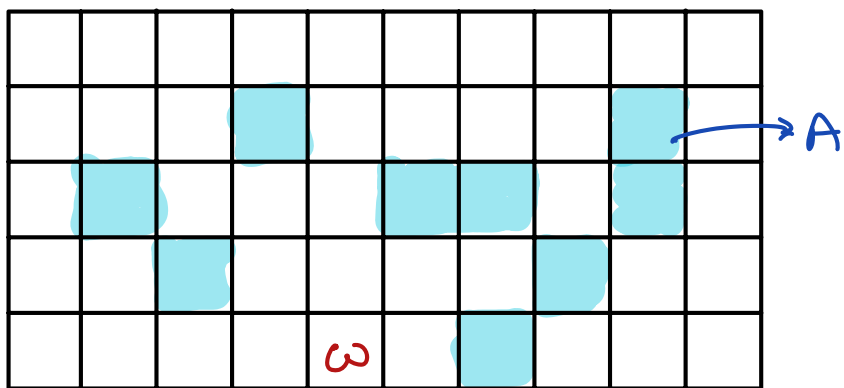


Why?

- Expectations (formally)
- Distributions
- ML / Information Theory

(Finite) Probability Spaces

Specified by a set Ω , and a "measure" $\nu: \Omega \rightarrow [0,1]$



$$\nu: \Omega \rightarrow [0,1]$$

$$\text{s.t. } \sum_{\omega \in \Omega} \nu(\omega) = 1$$

"probability measure"
"distribution"

Ω "sample space"
"outcome space"

Event: A subset of outcomes. $A \subseteq \Omega$

$$\mathbb{P}[A] = \sum_{\omega \in A} \nu(\omega)$$

$$\{\omega\} \subseteq \Omega$$

"atomic" or "elementary"
events

Random Variables (r.v.s)

Function on Ω

A (real-valued) random variable
is a **function**

$$X : \Omega \longrightarrow \mathbb{R}$$

Events $\equiv$ $\{0,1\}$ -valued r.v.s
indicator r.v.s

$$A \subseteq \Omega \equiv X_A(\omega) = \begin{cases} 1 & \omega \in A \\ 0 & \omega \notin A \end{cases}$$

Expectations

For probability space (Ω, ν) , r.v. $X: \Omega \rightarrow \mathbb{R}$

$$\mathbb{E}[X] = \sum_{\omega \in \Omega} \nu(\omega) X(\omega) = \sum_a \left(\underbrace{\sum_{\omega: X(\omega)=a} \nu(\omega)}_{P[X=a]} \right) \cdot a$$

Linearity of expectation

$$\mathbb{E}[X + Y] = \mathbb{E}[X] + \mathbb{E}[Y], \quad \mathbb{E}[c \cdot X] = c \cdot \mathbb{E}[X]$$

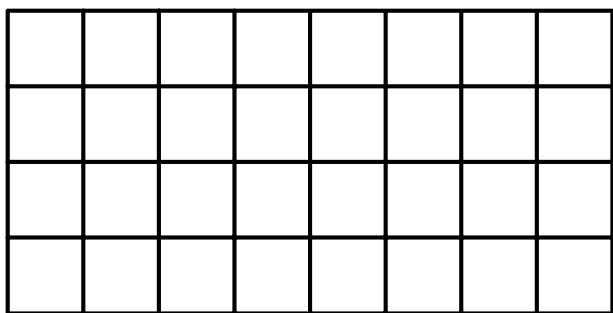
$V = \{X: \Omega \rightarrow \mathbb{R}\}$ is a vector space

$\mathbb{E}: V \rightarrow \mathbb{R}$ is a linear transformation

Shuffling Cards

Shuffle new deck of n cards to get a "random order"

How many cards in same place as before?



$X: \Omega \rightarrow \mathbb{R} = \# \text{ cards in same place.}$

$$E[X] = E[X_1 + \dots + X_n]$$

$\Omega = \text{Set of all } n! \text{ orderings}$

$X_i = i^{\text{th}} \text{ card in same place}$

$$P(\omega) = \frac{1}{n!}$$

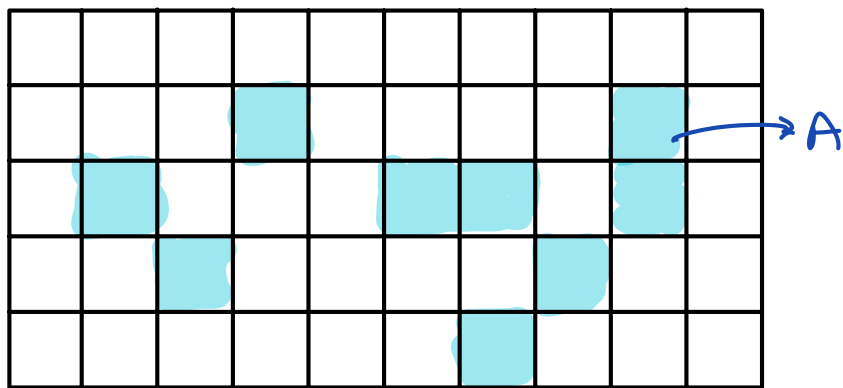
This is a common choice, but certainly not the only one!

$$E[X_i] = \frac{(n-1)!}{n!} = \frac{1}{n}$$

$$E[X] = 1$$

Conditioning

Conditioning on event $A \subseteq \Omega \equiv$ Restricting probability space to A



$$v_A(\omega) = \begin{cases} 0 & \text{if } \omega \notin A \\ \frac{v(\omega)}{P[A]} & \text{if } \omega \in A \end{cases}$$

$$P[B|A] = \sum_{\omega \in B} v_A(\omega)$$

$$= \sum_{\omega \in B} \frac{v(\omega)}{P[A]} \cdot \mathbb{1}_A(\omega)$$

$$= \sum_{\omega \in A \cap B} \frac{v(\omega)}{P[A]} = \frac{P[A \cap B]}{P[A]}$$

$$E[X|A] = \sum_{\omega} v_A(\omega) X(\omega)$$

$$= \frac{E[\mathbb{1}_A(\omega) \cdot X(\omega)]}{P[A]}$$